

Uniform Words are Primitive

Doron Puder

Department of Mathematics, Hebrew University

doronpuder@gmail.com

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Outline

- ▶ Uniform words, Primitive words

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- ▶ Some ingredients of the proof

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Definition

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- ▶ Primitive words are rare.

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w is primitive $\implies w$ is uniform

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- ▶ \exists many similar and extended open problems

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Theorem 1 (2011)

The conjecture holds for \mathbf{F}_2 .

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The conjecture holds for $\mathbf{F}_k \forall k$.

- ▶ *The proof involves: Stallings core graphs, random covering spaces, Möbius inversions, algebraic extensions of free groups,...*

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$$\begin{array}{c} w \text{ is primitive} \\ \Downarrow \textit{Observation} \\ \alpha_G(w) \sim U(G) \quad \forall \text{ finite } G \end{array}$$

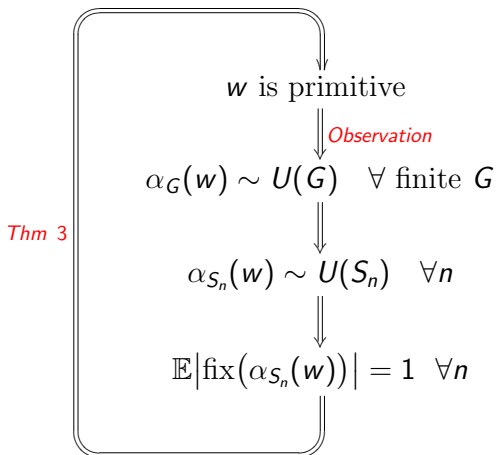
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E.g. $w = a^2b^2c^2$:

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3: the primitivity rank of w

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- ▶ w not primitive in $\mathbf{F}_k \implies \pi(w) \leq k$
- ▶ Thus, $\pi(w) \in \{0, 1, 2, \dots, k\} \cup \{\infty\}$

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$$\text{E.g. } \pi([a, b]) = \pi(a^2b^2) = 2$$

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$$\text{E.g. } \pi([a, b]) = \pi(a^2b^2) = 2$$

$$\pi(w_1w_2) = \pi(w_1) + \pi(w_2) \text{ for words with disjoint letters}$$

⋮

$$\pi(w) = \infty \iff w \text{ is primitive}$$

Primitivity Rank: a new classification of words

$$\pi(w) = \min \left\{ rk(J) \mid \begin{array}{l} w \in J \leq \mathbf{F}_k \text{ s.t.} \\ w \text{ is not primitive in } J \end{array} \right\}$$

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$\pi(w)$	Description	$\mathbb{E} \left[\begin{array}{l} \# \text{ fixed points} \\ \text{of } \alpha_{S_n}(w) \end{array} \right]$
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3. 2 new criteria (& algos) to detect primitivity (& free factors)

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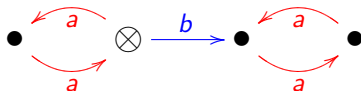
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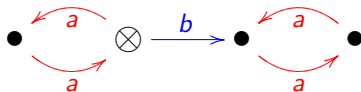
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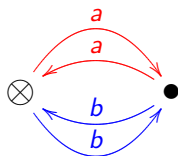
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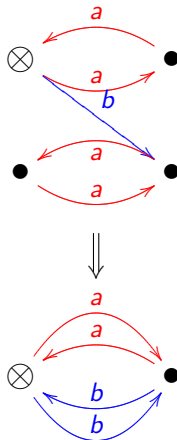
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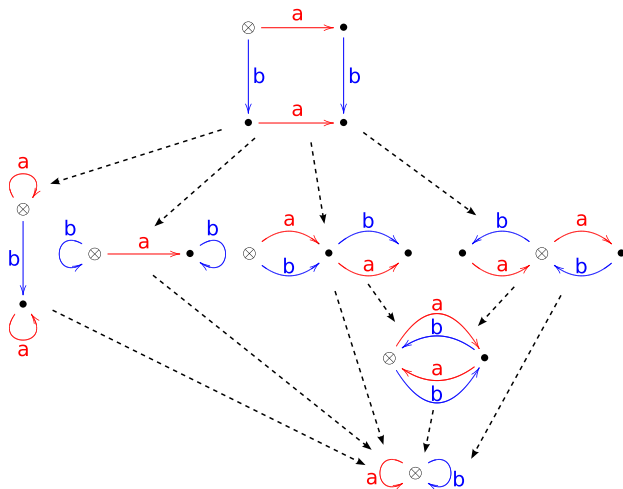
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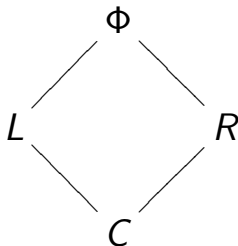
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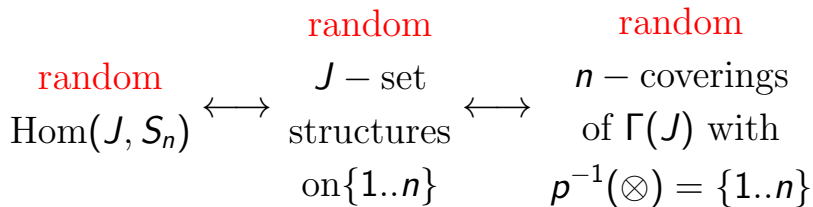
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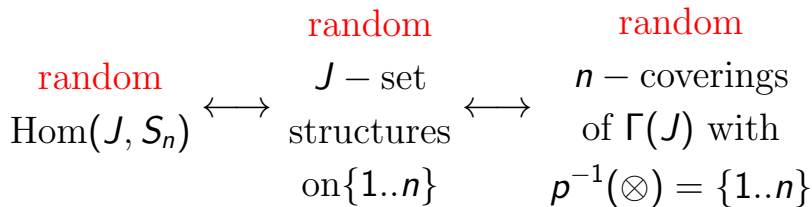
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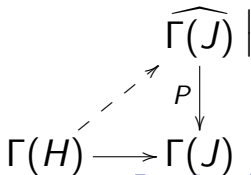
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Lemma: $\Phi_{H,J}(n) = \mathbb{E} \left| \text{lifts of } \Gamma(H) \text{ to } \widehat{\Gamma(J)} \right|$



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Thank You!